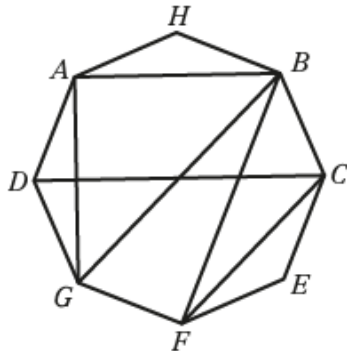
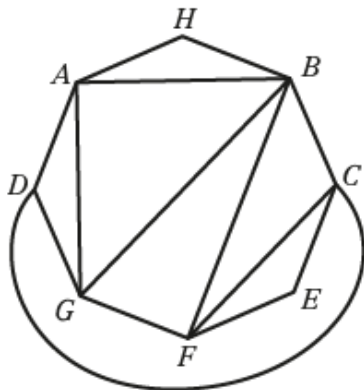


A Level Practice paper

- 1 a Starting with AB means that H cannot be included without repeating either B or A .
- b $AHBCEFGDA$
- c Draw a polygon matching the Hamiltonian cycle and look at the inside edges one at a time.



$AB(I), AG, BF, BG, CD, CF$
 $AB(I), AG(I), BF, BG, CD, CF$
 $AB(I), AG(I), BF(I), BG(I), CD(O), CF$
 $AB(I), AG(I), BF(I), BG(I), CD(O), CF(I)$
 The graph is planar.



2 a

42	31	36	18	27	33	41	47	12	24	16
42	36	41	47	33	31	18	27	12	24	16
42	47	41	36	33	31	18	27	24	16	12
47	42	41	36	33	31	27	18	24	16	12
47	42	41	36	33	31	27	24	18	16	12
47	42	41	36	33	31	27	24	18	16	12

All of the numbers have now been selected as pivots, so the list is in order.

- b** Using the first fit decreasing algorithm with reels of size 80

47 → 1

42 → 2

41 → 3

36 → 4

33 → 1 making it 80

31 → 2 making it 73

27 → 3 making it 68

24 → 4 making it 60

18 → 4 making it 78

16 → 5

12 → 5 making it 28

			18	
33	31	27	24	12
47	42	41	36	16

5 reels are required.

c $0.034 \times \frac{5000 \log 5000}{800 \log 800} = 0.27$ seconds

This is only an estimate because the time taken is only approximately proportional to $n \log n$

- 3 a** A network is semi-Eulerian if it has exactly two odd nodes.

- b** The odd nodes are A, B, C, D, E, F but we can ignore E and B .

Possible pairings: $AC + DF = 10 + 10 = 20$

$AD + CF = 11 + 9 = 20$

$AF + CD = 2 + 10 = 12$

So we use AF and CD

Minimum length = $72 + 12 = 84$ miles

- c** AF, CD

- d** This time we can ignore E and C

Possible pairings: $AB + DF = 4 + 10 = 14$

$AD + BF = 11 + 5 = 16$

$AF + BD = 2 + 13 = 15$

So we use AB and DF

The route is extended by $14 - 12 = 2$ miles

4 a Dijkstra's or Floyd's Algorithm

b Arc AC because of lack of symmetry in the table

c 1st iteration (no change)

-	5	∞	∞	∞
5	-	11	16	∞
10	11	-	8	6
∞	16	8	-	10
∞	∞	6	10	-

A	B	C	D	E
A	B	C	D	E
A	B	C	D	E
A	B	C	D	E
A	B	C	D	E

2nd iteration

-	5	16	21	∞
5	-	11	16	∞
10	11	-	8	6
21	16	8	-	10
∞	∞	6	10	-

A	B	B	B	E
A	B	C	D	E
A	B	C	D	E
B	B	C	D	E
A	B	C	D	E

3rd iteration

-	5	16	21	22
5	-	11	16	17
10	11	-	8	6
18	16	8	-	10
16	17	6	10	-

A	B	B	B	C
A	B	C	D	C
A	B	C	D	E
C	B	C	D	E
C	C	C	D	E

4th iteration (no change)

-	5	16	21	22
5	-	11	16	17
10	11	-	8	6
18	16	8	-	10
16	17	6	10	-

A	B	B	B	C
A	B	C	D	C
A	B	C	D	E
C	B	C	D	E
C	C	C	D	E

5th iteration (no change)

-	5	16	21	22
5	-	11	16	17
10	11	-	8	6
18	16	8	-	10
16	17	6	10	-

A	B	B	B	C
A	B	C	D	C
A	B	C	D	E
C	B	C	D	E
C	C	C	D	E

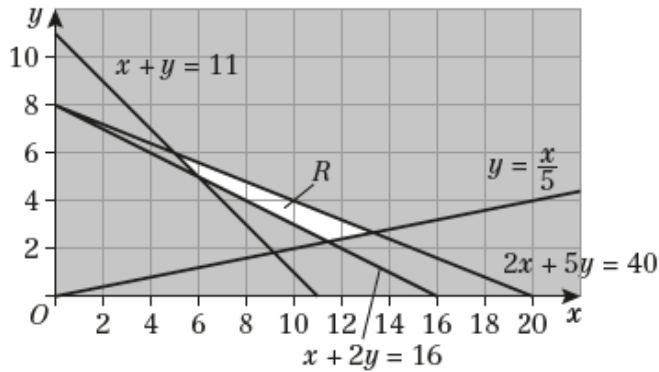
d In final iteration look at row A column E – goes through C

Now look at row A column C – goes through B

So $ABCE$ and length 22

5 a $12x + 30y \leq 240$, i.e. $2x + 5y \leq 40$
 $12x + 30y \leq 2(5x + 20y)$, i.e. $x \leq 5y$ or $y \geq \frac{x}{5}$
 $5x + 5y \geq 55$, i.e. $x + y \geq 11$
 $6x + 12y \geq 96$, i.e. $x + 2y \geq 16$

b

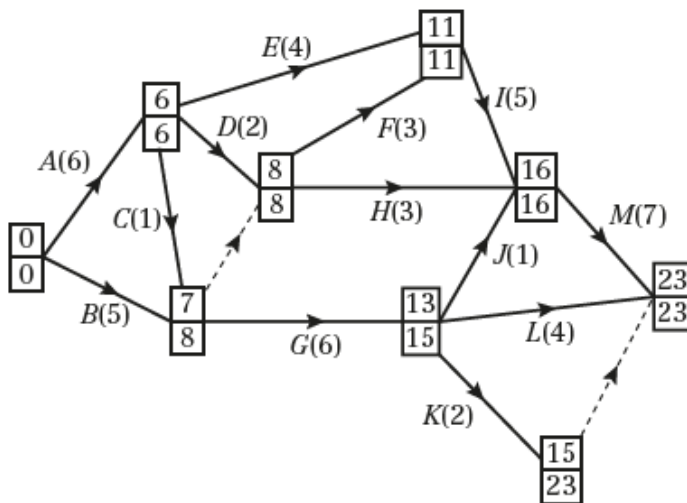


c

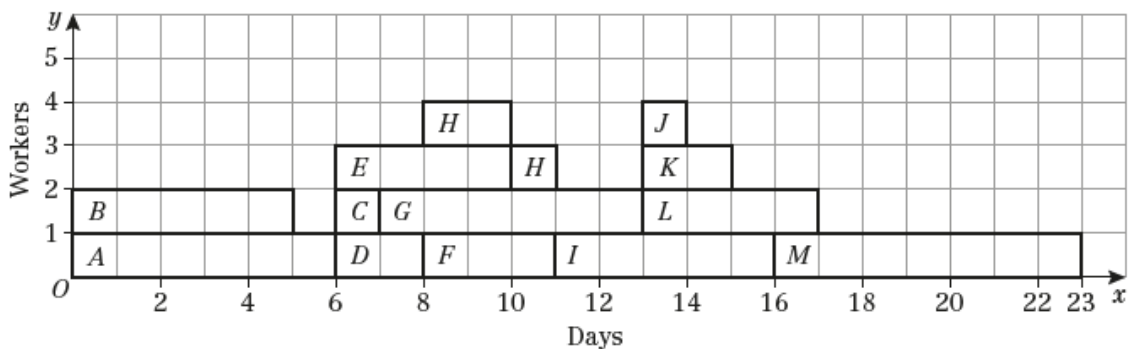
Vertex	(6, 5)	(5, 6)	$(\frac{80}{7}, \frac{16}{7})$	$(\frac{40}{3}, \frac{8}{3})$
Profit	£153	£166	£139.43	£162.67

So Angie should make 5 mini-packs and 6 mega-packs and the corresponding profit will be £166.

6 a



b



So 4 workers are required

c Delay the start of activity H by 2 days and the start of activity J by 2 days.

7 a The purpose of the first stage is to provide a basic feasible solution as a starting point for the second stage.

b Using slack, surplus and artificial variables we obtain:

$$2x + y + z + s_1 = 50$$

$$x + 3y + z + s_2 = 60$$

$$x - s_3 + a_1 = 10$$

This gives the first 3 lines of the tableau

$$P - x - 2y - z = 0 \text{ gives the 4}^{\text{th}} \text{ line}$$

We want to maximise $I = -a_1 = x - s_3 - 10$

$$\Rightarrow I - x + s_3 = -10$$

which gives the 5th line

b.v.	x	y	z	s ₁	s ₂	s ₃	a ₁	Value
s ₁	2	1	1	1	0	0	0	50
s ₂	1	3	1	0	1	0	0	60
a ₁	1	0	0	0	0	-1	1	10
P	-1	-2	-1	0	0	0	0	0
I	-1	0	0	0	0	1	0	-10

c So setting $s_3, y, z = 0$ we obtain the first solution of

$$I = 0 \text{ when } a_1 = 0, s_1 = 30, s_2 = 50$$

d x now replaces a_1 as a basic variable so we obtain the following table

b.v.	x	y	z	s ₁	s ₂	s ₃	Value
s ₁	0	1	1	1	0	2	30
s ₂	0	3	1	0	1	1	50
x	1	0	0	0	0	-1	10
P	0	-2	-1	0	0	-1	10

The smallest negative value in the bottom row is -2

So the y -column is the pivot column

The smallest value of θ is then $50 \div 3$

So the pivot is the 3 in the y -column

e The iterations are as follows:

bv	x	y	z	s ₁	s ₂	s ₃	Value
s ₁	0	1	1	1	0	2	30
s ₂	0	3	1	0	1	1	50
x	1	0	0	0	0	-1	10
P	0	-2	-1	0	0	-1	10

bv	x	y	z	s ₁	s ₂	s ₃	Value	Row operations
s ₁	0	1	1	1	0	2	30	
y	0	1	$\frac{1}{3}$	0	$\frac{1}{3}$	$\frac{1}{3}$	$\frac{50}{3}$	$R_2 \div 3$
x	1	0	0	0	0	-1	10	
P	0	-2	-1	0	0	-1	10	

bv	x	y	z	s ₁	s ₂	s ₃	Value	Row operations
s ₁	0	0	$\frac{2}{3}$	1	$-\frac{1}{3}$	$\frac{5}{3}$	$\frac{40}{3}$	$R_1 - R_2$
y	0	1	$\frac{1}{3}$	0	$\frac{1}{3}$	$\frac{1}{3}$	$\frac{50}{3}$	
x	1	0	0	0	0	-1	10	
P	0	0	$-\frac{1}{3}$	0	$\frac{2}{3}$	$-\frac{1}{3}$	$\frac{130}{3}$	$R_4 + 2R_2$

bv	x	y	z	s ₁	s ₂	s ₃	Value	Row operations
z	0	0	1	$\frac{3}{2}$	$-\frac{1}{2}$	$\frac{5}{2}$	20	$R_1 \times \frac{3}{2}$
y	0	1	$\frac{1}{3}$	0	$\frac{1}{3}$	$\frac{1}{3}$	$\frac{50}{3}$	
x	1	0	0	0	0	-1	10	
P	0	0	$-\frac{1}{3}$	0	$\frac{2}{3}$	$-\frac{1}{3}$	$\frac{130}{3}$	

bv	x	y	z	s ₁	s ₂	s ₃	Value	Row operations
z	0	0	1	$\frac{3}{2}$	$-\frac{1}{2}$	$\frac{5}{2}$	20	
y	0	1	0	$-\frac{1}{2}$	$\frac{1}{2}$	$-\frac{1}{2}$	10	$R_2 - \frac{1}{3}R_1$
x	1	0	0	0	0	-1	10	
P	0	0	0	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	50	$R_4 + \frac{1}{3}R_1$

So solution is $P = 50$ when $x = 10$, $y = 10$, $z = 20$