

## Route inspection 4A

1 a

vertex	A	B	C	D	E	F
valency	2	3	2	3	3	3

There are 4 nodes with odd valency so the graph is *neither* Eulerian nor Semi-Eulerian.

b

vertex	G	H	I	J	K
valency	3	4	3	2	4

There are precisely 2 nodes of odd degree (G and I) so the graph is *semi-Eulerian*.

A possible route starting at G and finishing at I is:

G – H – K – I – J – K – G – H – I

c

vertex	L	M	N	P	Q	R
valency	2	4	2	4	2	4

All vertices have even valency, so the graph is *Eulerian*.

A possible route starting and finishing at L is:

L – M – N – P – M – R – P – Q – R – L

2 a i

vertex	A	B	C	D	E	F	G	H
valency	4	2	4	2	2	4	2	2

ii

vertex	A	B	C	D	E	F	G
valency	4	4	2	4	2	4	4

b i A possible route is: A – B – C – A – F – C – E – G – H – F – D – A

ii A possible route is: A – C – F – A – B – E – G – B – D – G – F – D – A

3 a i

vertex	R	S	T	U	V	W
valency	2	2	3	3	2	2

Precisely 2 vertices of odd valency (T and U) so semi-Eulerian.

ii

vertex	H	I	J	K	L	M	N
valency	2	4	3	2	3	4	4

Precisely 2 nodes of odd degree (J and L) so semi-Eulerian.

b i A possible route starting of T and finishing at U is:

T – R – S – U – W – V – T – U

ii A possible route starts at J and finishes at L:

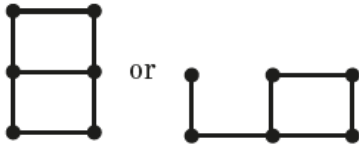
J – K – L – M – J – I – M – N – I – H – N – L

4 a The number of odd nodes of any graph must be even so this is not possible as there are 3 odd nodes.

b i  $2x + 1 + 2x + 4x - 1 + 4x + 6x = 2E = 18$   
 $\Rightarrow 18x = 18$   
 $\Rightarrow x = 1$

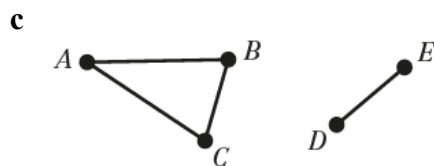
ii Semi-Eulerian since there are two odd nodes.

c Numerous possible answers e.g.:



5 a Not connected. There are no connections from  $A, B$  or  $C$  to  $D$  or  $E$ .

b Neither. To be Eulerian or semi-Eulerian the graph must be connected.



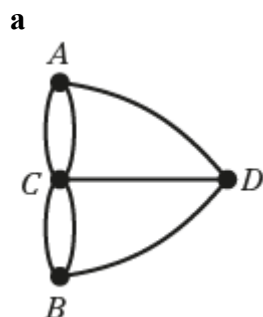
6 Adding up the numbers in each row, the orders of  $A, B, C, D, E$  are 2, 2, 2, 4, 4. Since they are all even the graph must be Eulerian.

7 a  $n$  must be odd so that each vertex will have degree  $n - 1$  which is even.



8 The example given in the question 1a is a counterexample.  $ABEFCDA$  is a Hamiltonian cycle, but the graph is not Eulerian.

**Challenge**

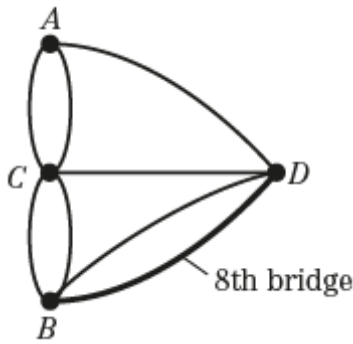


vertex	A	B	C	D
valency	3	3	5	3

There are more than two odd nodes, so the graph is *not* traversable.

Challenge

b

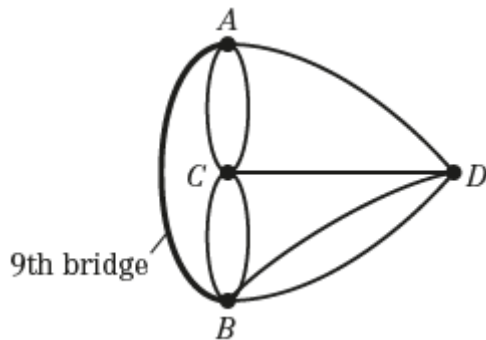


We will start at A and finish at C so these still need to have odd valency. We can only have two odd valencies so B and D must have even valencies (see table).

We need to change the valency of B and of D. So we build a bridge from B to D.

vertex	A	B	C	D
valency with 7 bridges	odd	<i>odd</i>	odd	<i>odd</i>
valency wanted	odd	<i>even</i>	odd	<i>even</i>

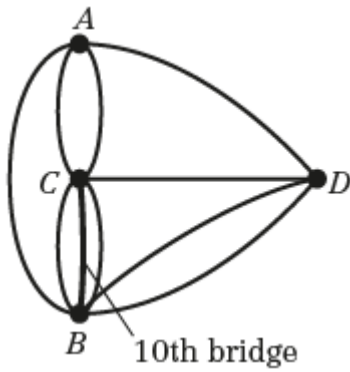
c



We will start at B and finish at C so these vertices need to be the two vertices with odd valency. We need A and D to have even valency (see table). We need to change the valency of node A and of node B. So we build a bridge from A to B.

vertex	A	B	C	D
valency with 8 bridges	<i>odd</i>	<i>even</i>	odd	even
valency wanted	<i>even</i>	<i>odd</i>	odd	even

Challenge  
d



All vertices now need to have even valency.  
 This means we need to change the valencies of nodes B and C.  
 So the 10th bridge needs to be built from B to C.

vertex	A	B	C	D
valency with 9 bridges	even	<i>odd</i>	<i>odd</i>	even
valency wanted	even	<i>even</i>	<i>even</i>	even